

Integer Computer Arithmetic

Computer Systems Sections 2.2, 2.3

Abstraction

Computer Integers behave like mathematical integers

Infinite & Ordered ($x < x+1$)

Associative $(x + y) + z = x + (y + z)$

Distributive $x*(y + z) = x*y + x*z$

Commutative $x+y = y+x, x*y=y*x$

Leaky Abstraction

Computer Integers behave like mathematical integers but...

Not infinite & ordered ($x < x+1$) e.g. (char) $127+1 = -128$

Associative $(x + y) + z = x + (y + z)$

Distributive $x*(y + z) = x*y + x*z$

Commutative $x+y = y+x$, $x*y=y*x$

Decimal, Binary, and Hexadecimal Bases

Dec	Bin	Hex
0	0b0000	0x0
1	0b0001	0x1
2	0b0010	0x2
3	0b0011	0x3
4	0b0100	0x4
5	0b0101	0x5
6	0b0110	0x6
7	0b0111	0x7

Dec	Bin	Hex
8	0b1000	0x8
9	0b1001	0x9
10	0b1010	0xA
11	0b1011	0xB
12	0b1100	0xC
13	0b1101	0xD
14	0b1110	0xE
15	0b1111	0xF

Non-Decimal Bases



$$value = \sum_{i=0}^{\#digits} d_i \times b^i, 0 \leq d_i < b$$

$$123 = 1 \times 10^2 + 2 \times 10^1 + 3 \times 10^0$$

$$0b0111\ 1011 = 2^6 + 2^5 + 2^4 + 2^3 + 2^1 + 2^0 = 64 + 32 + 16 + 8 + 2 + 1 = 123$$

$$0x7B = 7 \times 16 + 11 = 112 + 11 = 123$$

Notation

- A decimal number will be represented by decimal digits
 - $327 = 3 \times 100 + 2 \times 10 + 7 = 3 \times 10^2 + 2 \times 10^1 + 7 \times 10^0$
- A binary number will be prefixed by 0b, followed by 0/1
 - $0b0110 = 1 \times 4 + 1 \times 2 = 1 \times 2^2 + 1 \times 2^1 = 6$
- A hexadecimal number will be prefix by 0x, followed by hex digits
 - Hex digits are 0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F
 - $0x01AC = 1 \times 256 + 10 \times 16 + 12 = 1 \times 16^2 + 10 \times 16^1 + 12 \times 16^0 = 428$

Computer Unsigned Integer Arithmetic

- Represent a number as a vector of bits (binary data)

$$value = \sum_{i=0}^{n-1} B_i \times 2^i$$

B_{n-1}	B_{n-2}					B_1	B_0
0	0	0	1	0	1	0	1

$$0b0001\ 0101 = 0x15 = 1 \times 2^4 + 1 \times 2^2 + 1 \times 2^0 = 16 + 4 + 1 = 21$$

How Many Bits in an Integer in C?

Data Type	n	Max Value ²
char	8^1	UCHAR_MAX
short int	≥ 16	USHRT_MAX
int	$\geq 16^3$	UINT_MAX
long int	≥ 32	ULONG_MAX
long long int	≥ 64	ULLONG_MAX

¹ Typical value assuming addressable unit is 8 bits. Use “CHAR_BIT” in <limits.h> to get size of character.

² Constant defined when you “#include <limits.h>”

³ Usually 32

See http://en.wikipedia.org/wiki/C_data_types

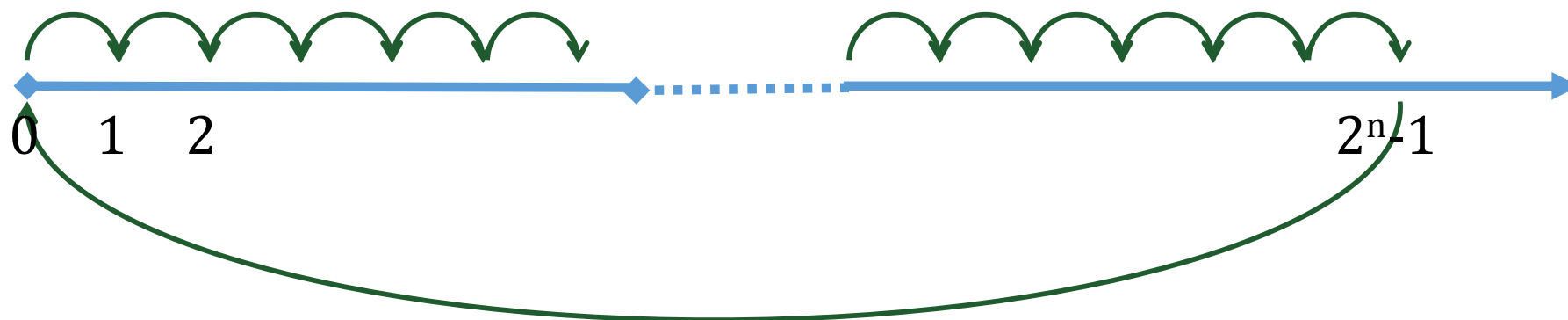
What happens on unsigned overflow?

- Overflowing bits ignored

$(\text{uchar})\ 0\text{xff} + (\text{uchar})\ 0\text{x01} = 0\text{x100} = (\text{uchar})\ 0\text{x00}$

$(\text{uchar})\ 255 + (\text{uchar})\ 1 = 256 = (\text{uchar})\ 0$

- See [xmp_counting/count.c](http://xmp.counting/count.c)

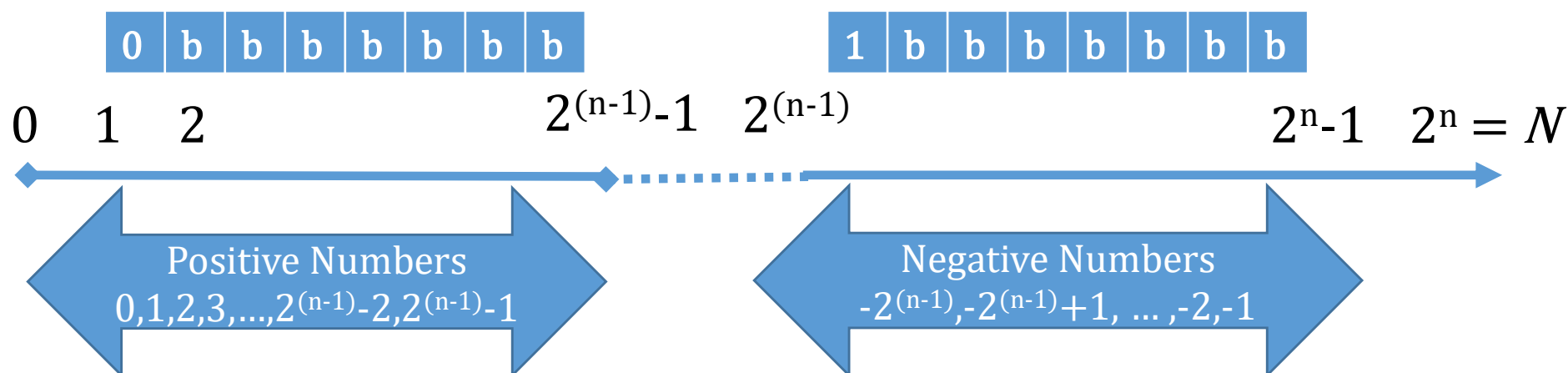


Unsigned: 8 bits

Hex	0x00	0x01	0x02	...	0x7E	0x7F	0x80	0x81	...	0xFE	0xFF
Binary	0000 0000	0000 0001	0000 0010		0111 1110	0111 1111	1000 0000	1000 0001		1111 1110	1111 1111
Unsigned Value	0	1	2	...	126	127	128	129	...	254	255

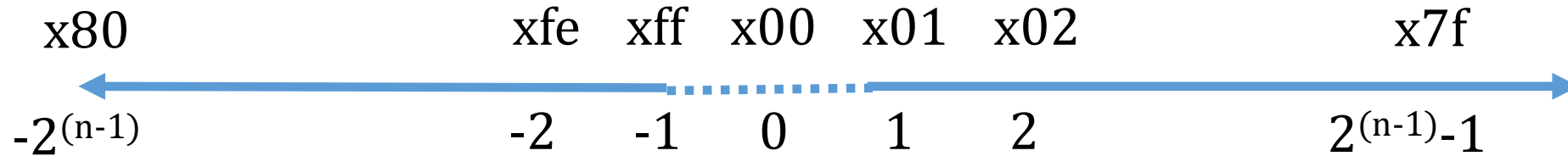
Computer Representation of Negatives

- Vector of 1's and 0's can't have "-"
- Trick : Subtract from N
 - N is bigger than any number you can represent
 - If you have n bits for a number, $N=2^n$
 - Given $x < 0$, represent x using $N - |x|$ or $N - (-x)$



Signed Integers

- Represented as “Two’s Complement” Numbers in n bits
- Positive numbers from 1 to $2^{(n-1)} - 1$: $valuebits = \sum_{i=0}^{n-2} B_i \times 2^i$
- Negative numbers from -1 to $-2^{(n-1)}$: $valuebits = 2^n - (\sum_{i=0}^{n-1} B_i \times 2^i)$
- Negative numbers have a “1” in the high order bit
- -1 is 0xffff...ff
- Shortcut... $-x =$ “Flip the bits and add 1”



Relationship between *bits* and *value*

- For $value \geq 0$
 - $bits = \sum_{i=0}^{n-1} B_i \times 2^i = value$
- For $value < 0$
 - $bits = 2^n - |value|$
 - $bits = 2^n - (-value) \quad -value = \sum_{i=0}^{n-1} B_i \times 2^i$
 - $bits = 2^n - \sum_{i=0}^{n-1} B_i \times 2^i$
 - $bits = (2^n - 1) - \left(\sum_{i=0}^{n-1} B_i \times 2^i\right) + 1$
 - $value = 2^n - bits$

Two's Complement : 8 bits $N=2^8 = 256$

Hex	0x80	0x81	...	0xFE	Value 0xFF	0x00	0x01	0x02	...	0x7E	0x7F
Binary	1000 0000	1000 0001	...	1111 1110	1111 1111	0000 0000	0000 0001	0000 0010	...	0111 1110	0111 1111
Unsigned Value	128	129	...	254	255	0	1	2	...	126	127
256-UVal	-128	-127	...	-2	-1						
Signed Value	-128	-127	...	-2	-1	0	1	2	...	126	127

Why does the shortcut work?

- ~~$valuebits = 2^n - (\sum_{i=0}^{n-1} B_i \times 2^i)$~~ (Assuming +) Flip the Bits
- $2^n - (-bits) = (\sum_{i=0}^{n-1} B_i \times 2^i)$
- $(2^n - 1) - (-value) + 1 = (\sum_{i=0}^{n-1} B_i \times 2^i)$
- So, for example, for signed char (8 bits) value of -3...

-value, flipped

	1	1	1	1	1	1	1	1
-	0	0	0	0	0	0	1	1
	1	1	1	1	1	1	0	0
+	0	0	0	0	0	0	0	1
	1	1	1	1	1	1	0	1

$2^n - 1$

-value

+1

What does “Flip the bits” do?

- Same as binary subtraction from a vector of all 1's

	1	1	1	1	1	1	1	1
-	0	0	1	0	0	0	1	1
	1	1	0	1	1	1	0	0

- Vector of all 1's = $(\sum_{i=0}^{n-1} 2^i) = 2^n - 1$
- So, “flip the bits” is really $2^n - 1 - value = 2^n - value - 1$

Why does the shortcut work?

- For $value \geq 0$, $bits = \sum_{i=0}^{n-1} B_i \times 2^i = value$
 - $-value = 2^n - \sum_{i=0}^{n-1} B_i \times 2^i = (2^n - 1) - bits + 1$
 - Flip the Bits
 - +1
- For $value < 0$, $bits = 2^n - \sum_{i=0}^{n-1} B_i \times 2^i$ where $-value = \sum_{i=0}^{n-1} B_i \times 2^i$
 - $-value = 2^n - 2^n + \sum_{i=0}^{n-1} B_i \times 2^i$
 - $-value = 2^n - (2^n - \sum_{i=0}^{n-1} B_i \times 2^i) = (2^n - 1) - bits + 1$
 - Flip the Bits
 - +1

How Many Bits in an Integer?

Data Type	n	Min Signed ²	Max Signed ²	Max Unsigned ²
char	8 ¹	SCHAR_MIN	SCHAR_MAX	UCHAR_MAX
short int	≥16	SHRT_MIN	SHRT_MAX	USHRT_MAX
int	≥16 ³	INT_MIN	INT_MAX	UINT_MAX
long int	≥32	LONG_MIN	LONG_MAX	ULONG_MAX
long long int	≥64	LLONG_MIN	LLONG_MAX	ULLONG_MAX

¹ Typical value assuming addressable unit is 8 bits. Use “CHAR_BIT” in <limits.h> to get size of character.

² Constant defined when you “#include <limits.h>”

³ Usually 16, often 32

See [xmp_limits/limits.c](http://xmp.limits/limits.c)

See http://en.wikipedia.org/wiki/C_data_types

What happens on signed overflow?

- Overflows into sign bit

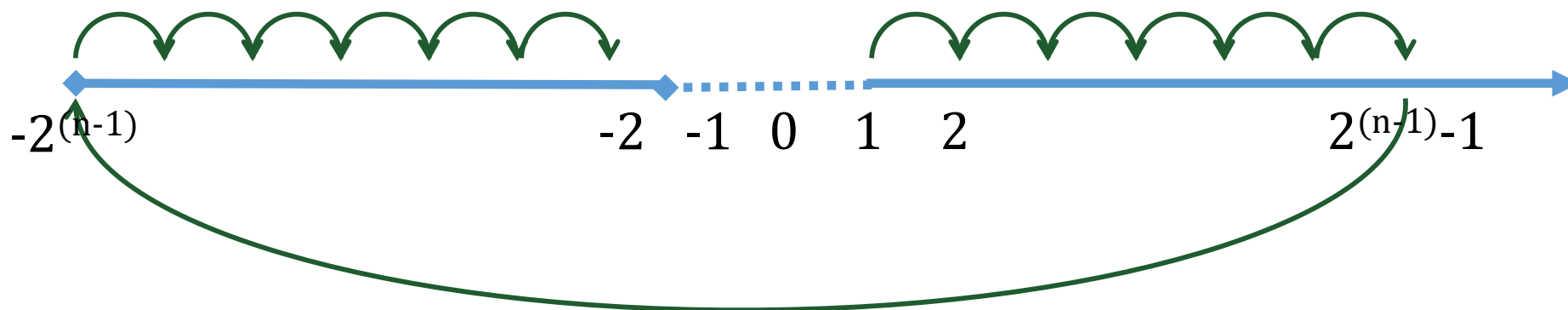
$(\text{char})\ 0x7f + (\text{char})\ 0x01 = 0x80 = (\text{char})\ 0x80$

$(\text{char})\ 127 + (\text{char})\ 1 = 128 = (\text{char})\ -128$

- See [xmp_counting/count.c](http://xmp.counting/count.c)

BITS

Value



Extending Positive Binary Numbers

- How do we make a number wider (more bits)?
- How do we add digits to a positive decimal number?
- 0023 (4 digits) -> ???????? (8 digits)
- 0b0101 (4 bits) -> 0b???? ???? (8 bits)

Extending Negative Binary Numbers

- ~~value~~*bits* = $2^n - \left(\sum_{i=0}^{n-1} B_i \times 2^i \right)$

$$\begin{array}{r}
 \text{n zeroes} \\
 2^n = \overbrace{1 \ 0 \ 0 \ 0 \ 0 \ \dots \ 0 \ 0 \ 0} \\
 - \quad \quad 0 \ 0 \ 1 \ 0 \ \dots \ 0 \ 1 \ 0 \\
 \hline
 1 \ 1 \ 0 \ 1 \ \dots \ 1 \ 0 \ 1
 \end{array}$$

$$\begin{array}{r}
 \text{m zeroes} \\
 2^m = \overbrace{1 \ 0 \ 0 \ 0 \ 0 \ \dots \ 0 \ 0 \ 0 \ \dots \ 0 \ 0 \ 0} \\
 - \quad \quad 0 \ 0 \ 0 \ 0 \ \dots \ 0 \ 1 \ 0 \ \dots \ 0 \ 1 \ 0 \\
 \hline
 1 \ 1 \ 1 \ 1 \ \dots \ 1 \ 0 \ 1 \ \dots \ 1 \ 0 \ 1
 \end{array}$$

Extending Binary Numbers

- Pad on the left with the sign bit

Truncating Binary Numbers

- Remove left most bits
- If any of the left most bits (including the resulting sign bit) are different from the sign bit, result is incorrect (OVERFLOW)
- Examples 8 bit to 4 bit...

0b 0000 0011 -> 0b 0011 (3->3)

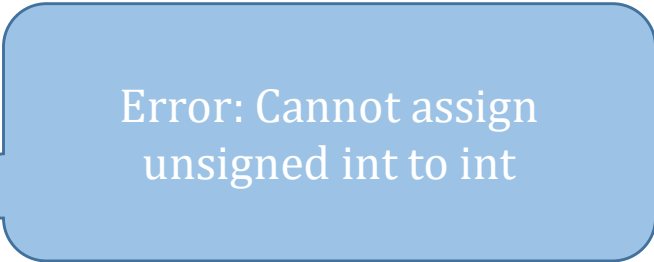
0b 0010 0011 -> 0b 0011 (35->3) OVFL

0b 1111 1001 -> 0b 1001 (-7 -> -7)

0b 1101 1001 -> 0b 1001 (-39 -> -7) OVFL

0b 1111 0111 -> 0b 0111 (-9 -> +7) OVFL

Mixing Types (“Casting”)

- How do you mix signed and unsigned numbers?
- How do you mix long / short numbers?
- Strong Typing – Compiler prevents mixing of types
`int x; unsigned int y; x=y;`
- Explicit Casting – Programmer tells compiler “treat this unsigned integer like a signed integer”
`int x; unsigned int y; x=(int) y;`
- Implicit Casting
`int x; unsigned int y; x=y;`

General Conversion Strategy

- If “from” size is smaller than “to” size (to precision is greater)
 - Pad on left with sign bit for signed data
 - Pad on left with 0’s for unsigned data
- If “from” size is greater than “to” size (to precision is less)
 - Truncate
- For signed->unsigned or unsigned->signed (same precision)
 - Bits stay the same... just interpreted differently.

Casting Examples

Legend

= precision

unsigned >precision

signed >precision

<precision

From			To			To			To		
Type	Bits	Value	Type	Bits	Value	Type	Bits	Value	Type	Bits	Value
uchar	0x01	1	char	0x01	1	uint	0x0001	1	int	0x0001	1
uchar	0xFF	255	char	0xFF	-1	uint	0x00FF	255	int	0xFFFF	-1
char	0x01	1	uchar	0x01	1	uint	0x0001	1	int	0x0001	1
char	0xFF	-1	uchar	0xFF	255	uint	0x00FF	255	int	0xFFFF	-1
uint	0xF0F3	61683	uchar	0xF3	243	char	0xF3	-13	int	0xF0F3	-3853
int	0x01D2	466	uchar	0xD2	210	char	0xD2	-46	uint	0x01D2	466

Integer Constant Types

- If not specified, integer constants are considered to be signed integers.
 - Make a constant “unsigned” by adding a “U” suffix
unsigned char x = 12U;

Implicit Casting for Integers in C

- In an expression,
 - if anything is unsigned, everything is unsigned!
`int x=-1; unsigned int y=321;`
`if (x>y) printf("I didn't expect this\n");`
 - everything smaller than int is cast to int
`char x=100; char y=50; x = (x*y)/500;`
 - everything smaller than the largest type is cast to the largest type
- In an assignment or function call,
 - the result is cast to the type of the receiver
- When in doubt, use explicit casts!

Abstraction

I don't need to worry about mixing types.
C takes care of implicit conversion
and does the right thing.

Leaky Abstraction

One of the most common causes of bugs in code occurs when naïve programmers assume C will do the right thing.

- Single unsigned variable causes all data to get treated as unsigned, especially in comparisons
- Assumption of floating point division when integer division (with truncation) is performed

Examples of implicit casting bugs

```
unsigned int length;  
int x;
```

```
if (x < length) {  
    /* what if x=-1, length=3 */  
    ...  
}
```

```
int grade;  
int test1, test2; // Scores <= 100  
  
grade = ((test1 + test2) / 200) * 100;  
// Everybody gets zero!
```

Conversion Error Summary

		TO TYPE							
		char	uchar	short	ushort	int	uint	long	ulong
FROM TYPE	char								
	uchar								
	short								
	ushort								
	int								
	uint								
	long								
	ulong								

LEGEND	
	Extend $\geq$ prec. No errors
	Extend $\geq$ prec. Wrong if <0
	Truncate $<$ prec Wrong if too big
	Truncate $<$ prec Wrong if <0 or too big